

Additional Verification Techniques

This material is from Chapter 4 in the textbook. The proof tableau scheme for *for-loops* is given in Section 4.1.

As an example we verify the partial correctness of the following:

```
ASSERT(0 <= n <= max)
{ int i;
  for (i=0; A[i] != x && i < n; i++)
    {}
  present = i<n;
}
ASSERT(present iff x in A[0:n-1])
```

Note that the code may not terminate normally if x does not occur in $A[0:n-1]$. Why?

As the loop invariant (denoted as I) we choose:

```
0<=i<=n && ForAll(k=0; k<i) x != A[k]
```

Using the scheme for for-loops, we must verify the following:

```
ASSERT(pre-condition)
i=0; /* initial assignment */
ASSERT(I) /*loop invariant*/
while( A[i] != x && i < n) {
  ASSERT(I && A[i] != x && i < n)
  /* for-loop has empty body*/
  i++;
  ASSERT(I)
}
ASSERT(I && !(A[i] != x && i < n))
```

```
present = i < n; /*final assignment*/
ASSERT(post-condition)
```

The complete construction is given in class. Here we have to be a little careful in how the post-condition is established from the invariant and the negation of the loop condition.

Array component assignment rule

The notation $(A \mid I \mapsto E)$ refers to an array obtained from A by replacing the value at position I by the value of the expression E .

More formally,

$$(A \mid I \mapsto E)[I'] = \begin{cases} E & \text{when } I' = I, \\ A[I'] & \text{when } I' \neq I. \end{cases}$$

Now the array component assignment rule can be written as:

$$[Q](A \mapsto A') \{A[I] = E;\} Q$$

where A' is $(A \mid I \mapsto E)$.

It has to be verified separately that the value of I is within the subscript range of the array A .

Example. We consider an array of even length.

The program should move elements from even numbered positions to a contiguous chunk at the beginning, see Figure 1.

The specification for the program is as follows:

```
Interface:  const int n;
            Entry A[2n]; /*entries numbered 0,...,2n-1 */
```

```
Pre-condition:  n >= 1 && A == A0
```

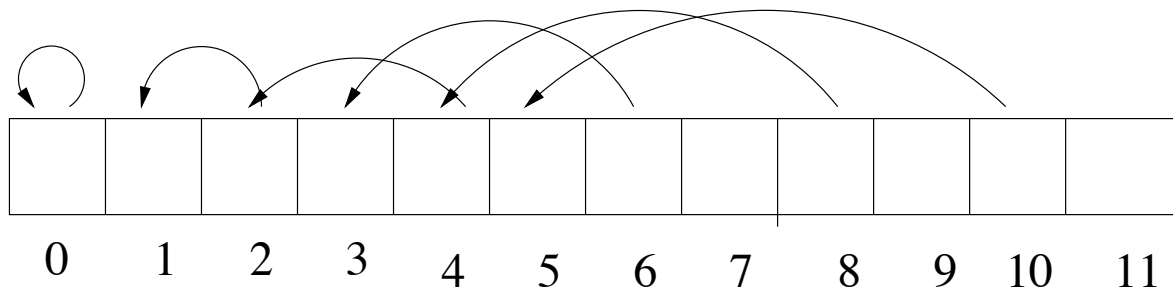


Figure 1: Moving of the array elements.

Post-condition: $\text{ForAll}(i=0; i<n) A[i] == A0[2i]$

On the basis of the post-condition we can select a suitable loop invariant and using it “derive” the program. (To be done in class.)